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C4-Sequences: Rate Adaptive Coded Modulation for Few Bits Message

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Outline





Binary Cyclic Code Shift Keying modulation

 $P_0 = 11101000 + BPSK modulation, q = 8, m = 3 bits per CCSK symbol.$

- CCSK modulation:
 - $P_0 = 11101000$
 - $P_1 = 01110100$
 - $P_2 = 00111010$
 - $P_3 = 00011101$
 - P₄ = 10001110
 - $P_5 = 01000111$
 - $P_6 = 10100011$
 - P₇ = 11010001

Code rate 3/8



[1] G.M. Dillard et al. « Cyclic code shift keying: a low probability of intercept communication technique." IEEE Trans. on Aerospace and Electronic Systems, 39(3):786–798, 2003.



More about CCSK

Element of the coding sequence are called "chip"

CCSK not restricted to BPSK modulation [1]: 8-PSK, 16-QAM, 64-APSK can be used to encode the chips.

Truncate the CCSK sequence to increase the code rate [1]:

Example: send only first I = 5 chips, instead of the whole sequence.

Binary message : 011001100Make 3-uplet symbols: $(011)_2(001)_2(100)_2$ Take decimal value: 3 1 4 Associate CCSK symbol P₃ P₁ P₄ Send => 000111010110010001110

• => 00011- - -01110- - -10001- - -

Code rate 3/5 • => 000110111010001

[1] G.M. Dillard et al. « Cyclic code shift keying: a low probability of intercept communication technique." IEEE Trans. on Aerospace and Electronic Systems, 39(3):786–798, 2003.

What is a good CCSK sequence?

- Notation: x = (x(0), x(1), ..., x(q-1)) a complex CCSK sequence with ||x||² = q (energy of 1 in average per chip).
- From **x** => define a codebook of *q* sequences $\mathbf{x}_a = (x(a) x(a+1), ..., x(a+q-1))$ (here, left rotation, with index operation done modulo q).

Minimum distance of the code: $D_m^2 = min_{a\neq b} || \mathbf{x}_a - \mathbf{x}_b ||^2$

$$||\mathbf{x}_{a} - \mathbf{x}_{b}||^{2} = ||\mathbf{x}_{a}||^{2} + ||\mathbf{x}_{b}||^{2} - 2 \operatorname{Real}(\langle \mathbf{x}_{a}, \mathbf{x}_{b} \rangle).$$

Thus : $D_m^2 = 2q - 2 \max(\text{Real}(< \mathbf{x}_a, \mathbf{x}_b>, a \neq b)))$



 How evolve the minimum distance when truncated sequences are used?

Convenient to use a normalized distance, i.e., a distance per chip, i.e.

$$D_m(l)^2 = \frac{1}{l} \min\left(\left\| \mathbf{x}_a^l - \mathbf{x}_b^l \right\|^2, a \neq b \right)$$

• With x_a^l the length-/truncated sequence of x_a , i.e., $x_a^l = (x(a), x(a + 1), ..., x(a + l - 1))$

Evolution of D_m(*I***) as a function of the truncation length?**



• It is possible to do better?



Outline

Summary of Marchand's presentation







Multi-objective optimization problem: maximize D_I² for I = 1, ..., q.



• Set of optimal solutions obtained.



Set of solutions with $D_m(I) = 2$ for several values of *I*.



Optimal for I = q/4, q/2, 3q/4 and q. Always above Binary and ZC sequences...



Outline

Summary of Marchand's presentation











Multi-objective optimization problem: maximize D_I² for I = 1, ..., q.



• Set of optimal solutions obtained.



What does C4 means?

o Constellation

• Circular





• (auto)-Correlation

q-1 $\sum x(n)x(n+\tau)'$ n=0

o Cross



Mathematical recipe for C4-sequence

Generation of a C4 sequence of length q = 4p in 3 steps.

1) Select randomly q/4 points on a circle of radius $r = 4\sqrt{q/4}$



2) From length q/4 vector **S**, generate length q vector **X** as

```
Function x = random_C4(q);
% Step 1
theta = 2*pi*rand(1,q/4);
S = 4*sqrt(q/4)*exp(1i*theta);
```

```
% Step 2
X = kron(S, [0 0 0 1]);
```

```
% Step 3
x = ifft(X);
```

 $\mathbf{X} = [0 \ 0 \ 0 \ S(0) \ 0 \ 0 \ S(1) \ \dots \ 0 \ 0 \ S(q/4-1)]$

3) Compute the Inverse Discrete Fourier Transform \mathbf{x} of \mathbf{X} , $\mathbf{x} = IDFT(\mathbf{X})$.

...and x is a C4 sequence



Experimentation Results q = 64







Clock-wise and Anti-clockwise C4 sequences

q = 64;	
theta1 = $rand(1,q/4)$;	theta2 = $rand(1,q/4)$;
S1 = 4*sqrt(q/4)* exp(2*pi*1i*theta1);	S2 = 4*sqrt(q/4)*exp(2*pi*1i*theta2);
X = kron(S1, [0 0 0 1]);	Y = kron(S2, [0 1 0 0]);
x = ifft(X);	y = ifft(Y);
Clock-wise C4 sequence	Anti-clockwise C4 sequence

Theorem: if x is a clockwise C4-sequence and y and anticlockwise C4-sequence, then,

for all a, b, $\langle x_a, y_b \rangle = 0$.

i.e., the two sets of sequences are orthogonal.



Properties of C4 sequences

• Invariant by a $\pi/2$ rotation (4-fold symmetry)

•
$$x\left(n+\frac{q}{4}\right) = j^{-c}x(n)$$
, c=1 clockwise,
c=-1, anticlockwise.

•
$$R_x(\tau) = j^{-kc}q$$
, when $\tau = \frac{kq}{4}$, 0 otherwise

•
$$D_l\left(k\frac{q}{4}\right) = 2, k = 1, 2, 3, 4$$

• From a length q C4 sequence, it is possible to generate a set of q lengths sequence of length $p = q/4 \{x_a^p\}, a = 0, 1, ..., q$ verifying $\|x_a^p - x_b^p\|^2 \ge q/2$ if $a \ne b$



Outline

Summary of Marchand's presentation



Geometrical shaping





- String art sequences are particular case of C4sequences.
- By the selection of vector θ , flexibility to shape the C4-sequence.
 - ♦ => points on unit circle
 - $\diamond =>$ points on 2 or more circles
 - ♦ => Optimization for a given criteria
 - Mutual information (MI) of the constellation alone in AWGN channel.



q = 64 C4-constellation, SNR of 10 dB



Capacity: 3.4594 bit/s/Hz

MI: 3.4192 bit/s/Hz

98.8379 % of capacity



q = 64 C4-constellation, SNR of 5 dB



Capacity: 2.0574 bit/s/Hz

MI: 2.0537 bit/s/Hz

99.8227 % of capacity



q = 64 C4-constellation, SNR of 0 dB



Capacity: 1.0000 bit/s/Hz

MI: 0.9998 bit/s/Hz

99.998 % of capacity



q=64 C4-constellation, SNR of -5 dB



Capacity: 0.396409 bit/s/Hz

MI: 0.396408 bit/s/Hz

99.999891 % of capacity

q=256 C4-constellation, SNR of 15 dB

Outline

Summary of Marchand's presentation

• Global spectral efficiency: $\frac{mR}{\bar{l}}$ bit/s/Hz.

Degree 4 Parity Check over GF(64)

Overall q = 64, m = 6 C4 sequence optimized for 5 dB of SNR spectral efficiency (bit/s/Hz) 10^{0} 9/2 = 4.5= 19/4 = 2.2518 24 10⁻¹ 3/2 = 1.5bits bits 9/8 = 1.12510⁻² 3/4 = 0.75FER Outer NB code: $l \rightarrow 2$ single parity over 3/8 = 0.37510⁻³ GF(64), R = 3/43/16 = 0.188l = 64, 60, 48 $l \Rightarrow 3$ Spectral efficiency: l = 32, 303/20 = 0.15 10^{-4} $l \neq 4$ l = 18, 16 $\frac{mR}{l} = \frac{6 \times 3/4}{l} = \frac{9}{2l}$ bit/s/Hz. 3/40 = 0.075l = 12, 10, 8= 610⁻⁵ -10 -5 0 5 10 15 -15 20 E_s/N_0

0.07 bit/s/Hz

Finite length simulation, q = 64, k = 120 bits, NB-LDPC R = 1/3.

• Finite length simulation confirmed theoretical results.

Potential applications of C4-sequences

- May be used for geometrical shaping
- May be used in multi-users environment to identify users.
- May be used, like ZC sequences, for synchronization.
- May be used for flexible rate adaptive coded modulation.

Potential extensions

• The 4 in C4 can also be the 4 optimal truncation lengths.

 Extension of C4 sequences to C3 or C5 sequences possible: in the construction process, replace Y = 4*sqrt(q/4)*kron(S, [0 1 0 0]);

Y = 5*sqrt(q/5)*kron(S, [0 1 0 0 0]); (for example) to create C5 sequence with a 5-fold symmetry

niv-ubs.fr/ https://ai4code.projects.labsticc.

Thank you !